Notes from 11/2

import Control.Monad.State

TBD: this needs work

Monads

A type is a monad m if it provides the following two operations:

```
return :: a -> m a
(>>=) :: m a -> (a -> m b) -> m b -- aka "bind"
returnSt a s = (a,s)
bindSt f g s0 = (b, s2)
 where (a, s1) = f s0
        (b, s2) = g a s1
Nothing >>= \_ -> Nothing
Just 14 >>= \ -> Nothing
Just 14 >>= a \rightarrow Just(a*2)
[4,5,6] >>= \a -> [a*2]
[4,5,6] >>= \a -> [a*2, a+5, 9]
xs >>= \a -> xs >>= \b -> xs >>= \c -> if c*c == a*a + b*b then [(a,b,c)] else []
"ABCD" >>= \c -> [succ c, pred c]
twiceM =
 do a <- [4,5,6]
     return (a*2)
pythag =
 do a <- [1..100]
     b <- [1..a]
     c <- [1..100]
     if c*c == a*a + b*b
     then return (a,b,c)
     else []
```

State monad

When we worked with threading state through a calculation or traversal, we defined a state function with a type like $s \rightarrow (r, s)$ where s is the state and r is the result. It turns out that this function type is *itself* a monad.

Substitute m a in the monad operations with $s \rightarrow (a, s)$ and you get:

return :: a -> s -> (a, s) (>>=) :: (s -> (a, s)) -> (a -> s -> (b, s)) -> s -> (b, s)

We can actually write these functions, but we'll name them a little differently so they don't conflict with the real ones:

Notice the usual careful state-passing in the where clause of the bindState function. That little bit of state-passing logic turns is sufficient to encode everything we need to build more elaborate stateful functions.

For example, the function threeRandoms becomes:

```
threeRandoms :: Seed -> ([Integer], Seed)
threeRandoms =
  rand `bindSt` \r1 ->
  rand `bindSt` \r2 ->
  rand `bindSt` \r3 ->
  returnState [r1, r2, r3]
```

This has the same type as the threeRandoms we did list week, but all the state-passing is hidden. It's done transparently by bindState and returnState, and we get exactly the same results.

```
ghci> threeRandoms (Seed 3453)
([58034571,429459059,225867046],Seed {unSeed = 225867046})
data Seed = Seed { unSeed :: Integer }
    deriving (Eq, Show)
```

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```
threeRandsSt :: State Seed [Integer]
threeRandsSt =
 do r1 <- randSt</pre>
    r2 <- randSt
    r3 <- randSt
     return [r1,r2,r3]
randSt :: State Seed Integer
randSt = do
  Seed s <- get
 let s' = (s * 16807) `mod` 0x7FFFFFF
 put (Seed s')
 return s'
rand :: Seed -> (Integer, Seed)
rand (Seed s) = (s', Seed s')
 where
    s' = (s * 16807) `mod` 0x7FFFFFF
main = putStrLn "OK"
```

"do" notation

Reader monad

```
Substitute r \rightarrow a form a.

returnRd :: a \rightarrow r \rightarrow a

returnRd a = a

bindRd :: (r \rightarrow a) \rightarrow (a \rightarrow r \rightarrow b) \rightarrow r \rightarrow b

bindRd f g r = g (f r) r

ask :: r \rightarrow r

ask r = r
```

r Represents an *environment* – a value that is always available to you, but you don't have to explicitly pass around.

Reader example: do some calculations on a list. Then take modulo N, where N is from the environment.

calc :: Integer -> Integer -> Integer
calc x = blop x `bindRd` \y -> returnRd (sum y)

```
blop :: Integer -> Integer -> [Integer]
blop x =
    cran x `bindRd` \y ->
    ask `bindRd` \n ->
    returnRd $ map (`mod` n) y
cran :: Integer -> Integer -> [Integer]
cran x = returnRd $ map (2^) [1..x]
```

Writer monad

Substitute (a,w) for m a.

```
returnW :: Monoid w => a -> (a,w)
returnW a = (a, mempty)
bindW :: Monoid w => (a,w) -> (a -> (b,w)) -> (b,w)
bindW (a,w1) f = (b, mappend w1 w2)
where (b,w2) = f a
tell :: Monoid w => w -> ((),w)
tell w = ((), w)
addEvens xs = sum $ filter even xs
```

I'd like this function to log the odd numbers that it skips when filtering.

```
filterW :: (a -> Bool) -> [a] -> ([a],[a])
filterW f [] = returnW []
filterW f (x:xs) =
  filterW f xs `bindW` \ys ->
  if f x
  then returnW (x:ys)
  else tell [x] `bindW` \() ->
     returnW ys
addEvensW xs =
  filterW even xs `bindW` \ys -> returnW (sum ys)
```